

Neighbour-Sum-2-Distinguishing Edge-Weightings: Doubling the 1-2-3 Conjecture[☆]

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Abstract

The 1-2-3 Conjecture asks whether every graph with no connected component isomorphic to K_2 can be 3-edge-weighted so that every two adjacent vertices u and v can be distinguished via the sum of their incident weights, that is the incident sums of u and v differ by at least 1.

We here investigate the consequences on the 1-2-3 Conjecture of requiring a stronger distinction condition. Namely, we consider two adjacent vertices distinguished when their incident sums differ by at least 2. As a guiding line, we conjecture that every graph with no connected component isomorphic to K_2 admits a 5-edge-weighting permitting to distinguish the adjacent vertices in this stronger way.

We prove this conjecture for several classes of graphs, including bipartite graphs and cubic graphs. We then consider algorithmic aspects, and show that it is NP-complete to determine the smallest k such that a given bipartite graph admits such a k -edge-weighting. In contrast, we show that the same problem can be solved in polynomial time for a given tree.

Keywords: 1-2-3 Conjecture; Difference-2 distinction; Bipartite graphs.

1. Introduction

Let G be a graph, and ω be an edge-weighting of G . For every vertex v , one can compute its *incident sum* $\sigma_\omega(v)$ (or simply $\sigma(v)$ when no ambiguity is possible) of weights by ω , being $\sigma(v) := \sum_{u \in N(v)} \omega(vu)$, where $N(v)$ denotes the set of neighbours of v . We call ω *neighbour-sum-distinguishing* if it yields a proper σ , *i.e.*, we have $\sigma(u) \neq \sigma(v)$ for every edge uv of G . It can be observed that every connected graph different from K_2 admits a neighbour-sum-distinguishing edge-weighting. Graphs with no connected component isomorphic to K_2 are thus said *nice*, with respect to neighbour-sum-distinguishing edge-weightings. For a nice graph G , it thus makes sense to investigate the smallest k such that G admits a neighbour-sum-distinguishing k -edge-weighting. This smallest k is denoted by $\chi_\Sigma(G)$.

The 1-2-3 Conjecture, addressed by Karoński, Łuczak and Thomason [4], asks whether $\chi_\Sigma(G) \leq 3$ holds for every nice graph G .

1-2-3 Conjecture. *For every nice graph G , we have $\chi_\Sigma(G) \leq 3$.*

If true, the bound in the 1-2-3 Conjecture would be best possible, as attested for example by nice complete graphs and cycles with length not multiple of 4. More generally,

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it is NP-complete to decide whether $\chi_\Sigma(G) \leq 2$ holds for a given graph G , as first proved by Dudek and Wajc [1]. The same problem, however, can be handled in polynomial time when restricted to bipartite graphs, as recently shown by Thomassen, Wu and Zhang [7]. To date, the best result towards the 1-2-3 Conjecture is due to Kalkowski, Karoński and Pfender [3], who proved that $\chi_\Sigma(G) \leq 5$ holds for every nice graph G .

When designing neighbour-sum-distinguishing edge-weightings, the goal is to make adjacent vertices distinguishable via their incident sums. In ordinary neighbour-sum-distinguishing edge-weightings, adjacent vertices are considered distinguished as soon as their incident sums are distinct. We here investigate edge-weightings that permit to distinguish the adjacent vertices in a stronger way. Namely, we require adjacent vertices to have incident sums differing by at least 2. An edge-weighting with this stronger requirement is said to be *neighbour-sum-2-distinguishing* throughout. As observed in upcoming Observation 2.1, a neighbour-sum-distinguishing k -edge-weighting can easily be turned into a neighbour-sum-2-distinguishing $2k$ -edge-weighting. Moreover, since K_2 does clearly not admit any neighbour-sum-2-distinguishing edge-weighting, the notion of nice graphs for neighbour-sum-distinguishing edge-weightings and for neighbour-sum-2-distinguishing edge-weightings coincide. Again, we can thus wonder about the smallest k such that a given nice graph G admits a neighbour-sum-2-distinguishing k -edge-weighting, which we denote by $\chi_{\Sigma>1}(G)$.

Our main goal in this paper is to study how $\chi_{\Sigma>1}$ behaves in general, in particular for graphs for which the parameter χ_Σ is well understood. As noted in upcoming Observation 2.1, the 1-2-3 Conjecture, if true, would imply that $\chi_{\Sigma>1}(G) \leq 6$ holds for every nice graph G . One could thus naturally wonder about a 1-2-3-4-5-6 Conjecture for neighbour-sum-2-distinguishing edge-weightings. It actually turns out that we did not manage to exhibit nice graphs G with $\chi_{\Sigma>1}(G) = 6$. On the other hand, we prove, throughout this paper, that for several common classes of nice graphs G we have $\chi_{\Sigma>1}(G) \leq 5$. We are thus tempted to address the following.

Conjecture 1.1. *For every nice graph G , we have $\chi_{\Sigma>1}(G) \leq 5$.*

We here give first evidence towards Conjecture 1.1. We start in Section 2 by raising connections between neighbour-sum-distinguishing edge-weightings and neighbour-sum-2-distinguishing edge-weightings, from which we deduce first bounds on $\chi_{\Sigma>1}$. In Section 3, we then prove Conjecture 1.1 for nice bipartite graphs, and investigate further some particular classes of bipartite graphs, from which we exhibit interesting aspects of the problem. The algorithmic aspects are considered in Section 4, where we first prove that it is NP-complete to determine the exact value of $\chi_{\Sigma>1}$. This statement is showed to remain true even for bipartite graphs, which contrasts with the complexity of determining the exact value of χ_Σ for these graphs. We then show that determining the exact value of $\chi_{\Sigma>1}$ can be done in polynomial time for trees. Perspectives for future works are gathered in Section 5.

2. Preliminaries

In this section, we point out general properties that actually apply to neighbour-sum- d -distinguishing edge-weightings (for any $d \geq 1$), *i.e.*, edge-weightings where the adjacent sums differ by at least d . For a given graph G , we denote by $\chi_{\Sigma>d-1}(G)$ the least integer k such that G admits a neighbour-sum- d -distinguishing k -edge-weighting. Note that the parameter $\chi_{\Sigma>0}$ actually corresponds to χ_Σ .

Obviously, by multiplying all weights assigned by a neighbour-sum-distinguishing k -edge-weighting by d , we obtain a neighbour-sum- d -distinguishing dk -edge-weighting.

Observation 2.1. *For every nice graph G and integer $d \geq 1$, we have $\chi_{\Sigma>d-1}(G) \leq d\chi_{\Sigma}(G)$. In particular, $\chi_{\Sigma>1}(G) \leq 2\chi_{\Sigma}(G)$.*

Observation 2.1 already has several implications towards Conjecture 1.1. First, the 1-2-3 Conjecture, if true, would imply that $\chi_{\Sigma>1}(G) \leq 6$ holds for every nice graph G . Although we still do not know whether the 1-2-3 Conjecture is true, every partial result towards that conjecture can be adapted to Conjecture 1.1. In that line, perhaps the most interesting result to consider is the one due to Kalkowski, Karoński and Pfender [3], who proved that $\chi_{\Sigma}(G) \leq 5$ holds for every nice graph G . In our context, this and Observation 2.1 yield the following, which shows that, when requiring adjacent sums to differ by at least d , the number of needed weights grows linearly with d .

Corollary 2.2. *For every nice graph G and integer $d \geq 1$, we have $\chi_{\Sigma>d-1}(G) \leq 5d$. In particular, $\chi_{\Sigma>1}(G) \leq 10$.*

The bound in Corollary 2.2 is immediately improved for every graph G for which we know that $\chi_{\Sigma}(G) < 5$ holds. In particular, we have $\chi_{\Sigma>1}(G) \leq 6$ for every nice graph G verifying the 1-2-3 Conjecture, which is very close to Conjecture 1.1. Let us recall, in particular, that the 1-2-3 Conjecture was verified for nice bipartite graphs, 3-chromatic graphs, nice complete graphs, and regular graphs with sufficiently large degree. We here refer the reader to the survey [6] by Seamone, wherein all such results are gathered.

As stated in Observation 2.1, by multiplying all weights assigned by a neighbour-sum-distinguishing k -edge-weighting by a same integer d , we get a neighbour-sum- d -distinguishing dk -edge-weighting since each $\sigma(v)$ is multiplied by d . This of course does not have to be true if one decreases (or increases) all weights by a same d , since, here, the effect on each $\sigma(v)$ depends on $d(v)$. There are situations, however, where this can be done safely.

Observation 2.3. *Let ω be a neighbour-sum-distinguishing d -edge-weighting, $d \geq 1$, of a graph G such that, for every edge uv (where $d(u) \geq d(v)$), we have $\sigma(u) < \sigma(v)$ (resp. $\sigma(u) > \sigma(v)$). Then, by decreasing (resp. increasing) all edge weights by a same integer x , we get another neighbour-sum-distinguishing d -edge-weighting of G .*

Due to the fact that, in the context of Conjecture 1.1, we focus on edge-weightings assigning strictly positive weights, when decreasing edge weights we should also make sure that none becomes null or negative. Observation 2.3 can nevertheless be used when the smallest edge weight value assigned by ω is known. As an illustration, we improve Observation 2.1 for nice regular graphs.

Corollary 2.4. *For every nice regular graph G and integer $d \geq 1$, we have $\chi_{\Sigma>d-1}(G) \leq d\chi_{\Sigma}(G) - (d - 1)$. In particular, $\chi_{\Sigma>1}(G) \leq 2\chi_{\Sigma}(G) - 1$.*

Corollary 2.4 notably implies that Conjecture 1.1 holds for nice complete graphs and 3-colourable regular graphs, as the 1-2-3 Conjecture was proved for such graphs. Other bounds also follow for regular graphs with larger degree, see [6].

3. Conjecture 1.1 for bipartite graphs

Since the 1-2-3 Conjecture holds for nice bipartite graphs (see [4]), from Observation 2.1 we directly get that $\chi_{\Sigma>1}(G) \leq 6$ holds for every nice bipartite graph G . It was actually

proved by Thomassen, Wu and Zhang [7] that a bipartite graph G satisfies $\chi_\Sigma(G) = 3$ if and only if G is an odd multicactus. *Odd multicacti* can be defined as follows. Start from a collection C_1, \dots, C_m of $m \geq 1$ cycles whose lengths are at least 6 and congruent to 2 modulo 4, and colour the edges of the C_i 's in a proper way using colours red and green. An odd multicactus is then any connected graph obtained by repeatedly applying the following operation: pick two connected components G_1 and G_2 , and identify a green edge of G_1 with a green edge of G_2 . Note that, in particular, every cycle whose length is congruent to 2 modulo 4 is an odd multicactus.

From this all, we directly get the following.

Corollary 3.1. *For every nice bipartite graph G , we have*

$$\chi_{\Sigma>1}(G) \leq \begin{cases} 4 & \text{if } G \text{ is not an odd multicactus,} \\ 6 & \text{otherwise.} \end{cases}$$

In the next result, we improve Corollary 3.1 by completely proving Conjecture 1.1 for all nice bipartite graphs.

Theorem 3.2. *For every nice bipartite graph G , we have $\chi_{\Sigma>1}(G) \leq 5$.*

Proof. Let A and B denote the two partite sets of G . We prove a stronger statement, namely that G admits a 5-edge-weighting such that the vertices from one of the partite sets have incident sum congruent to 2 or 3 modulo 5, while the vertices from the other partite set have incident sum congruent to 0 modulo 5. Such an edge-weighting is clearly neighbour-sum-2-distinguishing. In this setting, note that we can equivalently look for a neighbour-sum-2-distinguishing $\{0, 1, 2, 3, 4\}$ -edge-weighting.

Assume the edges of G are weighted in some way. In what follows, we repeatedly apply the following modification procedure for some given $\alpha \in \{1, 2, 3, 4\}$. Let u and v be two distinct vertices from a same partite set, and let P be a path (obviously of even length) from u to v in G . The procedure consists in modifying the weights of the edges of P from one end to the other by alternatively subtracting and adding α (modulo 5), *i.e.*, apply $-\alpha, +\alpha, -\alpha, \dots, -\alpha, +\alpha$ to the weights along P . Note that this only alters the incident sums modulo 5 of u and v , the two ends of P , which change by $-\alpha$ and $+\alpha$, respectively. Also, note that the “orientation” of P is important, as the sum of its first vertex is altered by $-\alpha$ while the sum of its last vertex is altered by $+\alpha$. Hence, in what follows, though G is not oriented, we consider paths *from* a vertex *to* another one, so that the first and last vertices are clearly identified.

We proceed as follows to obtain the claimed edge-weighting of G . Assume first that one of the two partite sets, say, A , has even size. Start from all edges being weighted 0. Then repeatedly consider two new vertices u and v of A , and apply the modification procedure above with $\alpha = 2$ onto a path from u to v in G . As stated, only the incident sums of u and v modulo 5 are altered, from 0 to 3 and 2, respectively. Once the process has been applied for all pairs of vertices, we get the desired neighbour-sum-2-distinguishing $\{0, 1, 2, 3, 4\}$ -edge-weighting of G .

Assume now that A and B both have odd size. Repeating the exact same arguments as in the previous case on pairs of vertices of, say, A , we can deduce a $\{0, 1, 2, 3, 4\}$ -edge-weighting of G where all vertices of A but one vertex u have incident sum congruent to 2 or 3 modulo 5, while all vertices of B have incident sum congruent to 0 modulo 5. In particular, also u has incident sum congruent to 0 modulo 5. If A has two vertices v_1, v_2 with incident sum congruent to 2 (resp. 3) modulo 5, then we consider one path P_1 from u to v_1 and one path P_2 from u to v_2 , and apply the modification procedure for $\alpha = 1$ (resp.

$\alpha = 4 = -1$ modulo 5) onto P_1 and P_2 . This affects the incident sums of v_1 and v_2 which are now congruent to 3 (resp. 2) modulo 5, while u now has incident sum congruent to 2 (resp. 3) modulo 5. The edge-weighting is thus as desired.

The only situation where the previous case does not apply is when $|A| = |B| = 3$. In that case though, it can easily be checked by hand that G always admits a neighbour-sum-2-distinguishing $\{0, 1, 2, 3, 4\}$ -edge-weighting. \square

Although Theorem 3.2 completely proves Conjecture 1.1 for nice bipartite graphs, in the rest of this section we study bipartite graphs further. Namely, we deduce the exact value of $\chi_{\Sigma>1}$ for nice paths and cycles, and prove a refinement of Conjecture 1.1 for odd multicacti. These results support some observations to be raised in concluding Section 5.

In all upcoming proofs, we will deal with adjacent degree-2 vertices; for such a configuration, the following observation obviously applies.

Observation 3.3. *Let ω be a neighbour-sum-2-distinguishing edge-weighting of a graph G , and uv be an edge such that $d(u) = d(v) = 2$. Assuming u' (resp. v') denotes the neighbour of u (resp. v) different from v (resp. u), the weights $\omega(u'u)$ and $\omega(v'v)$ differ by at least 2.*

3.1. Paths

We denote by P_ℓ the path of length ℓ . Therefore, the path $P_1 = K_2$ is not nice. In the next result, we determine the value of $\chi_{\Sigma>1}(P_\ell)$ for every $\ell \geq 2$.

Theorem 3.4. *For every path P_ℓ , $\ell \geq 2$, we have*

$$\chi_{\Sigma>1}(P_\ell) = \begin{cases} 2 & \text{if } \ell = 2, \\ 3 & \text{if } \ell > 2 \text{ and } \ell \equiv 0, 2, 3 \pmod{4}, \\ 4 & \text{otherwise.} \end{cases}$$

Proof. Recall that $\chi_{\Sigma>1}(P_\ell) \leq 4$ holds for every $\ell \geq 2$, by Corollary 3.1. Moreover, since $\chi_{\Sigma>1}(G) = 1$ if and only if G is a graph such that the degrees of every two adjacent vertices differ by at least 2, we get $\chi_{\Sigma>1}(P_\ell) \geq 2$ for every $\ell \geq 2$.

Let v_0, \dots, v_ℓ denote the vertices of the path P_ℓ , with $v_i v_{i+1}$ being an edge for every i , $0 \leq i \leq \ell - 1$. We clearly have $\chi_{\Sigma>1}(P_2) = 2$ since the weighting ω given by $\omega(v_0 v_1) = \omega(v_1 v_2) = 2$ is neighbour-sum-2-distinguishing. Suppose now that $\ell \geq 3$. Then P_ℓ has two adjacent vertices v_i and v_{i+1} with degree 2, and Observation 3.3 applies. For any 2-edge-weighting ω of P_ℓ , we have $\sigma(v_i) = \omega(v_{i-1} v_i) + \omega(v_i v_{i+1})$ and $\sigma(v_{i+1}) = \omega(v_i v_{i+1}) + \omega(v_{i+1} v_{i+2})$. Since $\omega(v_{i-1} v_i), \omega(v_{i+1} v_{i+2}) \in \{1, 2\}$, necessarily $\sigma(v_i)$ and $\sigma(v_{i+1})$ differ by at most 1, so that ω cannot be neighbour-sum-2-distinguishing.

If $\ell = 3$, then assigning successive edge weights 1, 3, 3 to the edges of P_3 is neighbour-sum-2-distinguishing as it yields successive incident sums 1, 4, 6, 3. So we may suppose from now on that $\ell \geq 4$. Under that assumption, note that a neighbour-sum-2-distinguishing 3-edge-weighting ω of P_ℓ cannot assign weight 2 since Observation 3.3 applies. We thus restrict our attention to $\{1, 3\}$ -edge-weightings of P_ℓ . Note that the edge $v_1 v_2$ (and similarly $v_{\ell-1} v_\ell$) must be assigned weight 3 so that $\sigma(v_0)$ and $\sigma(v_1)$ differ by at least 2. From this, observe that all neighbour-sum-2-distinguishing $\{1, 3\}$ -edge-weightings of P_ℓ are (up to directional symmetry) 1, 3, 3, 1, 1, 3, 3, \dots , 3, 3, 1, and 1, 3, 3, 1, 1, \dots , 3, 3, and 3, 3, 1, 1, 3, 3, 1, 1, \dots , 3, 3. The claim then follows from the fact that these edge-weightings only apply for particular values of ℓ modulo 4. \square

3.2. Cycles

For every $\ell \geq 3$, we denote by C_ℓ the cycle of length ℓ . In the next result, we determine the value of $\chi_{\Sigma>1}(C_\ell)$ for every cycle C_ℓ .

Theorem 3.5. *For every cycle C_ℓ , $\ell \geq 3$, we have*

$$\chi_{\Sigma>1}(C_\ell) = \begin{cases} 3 & \text{if } \ell \equiv 0 \pmod{4}, \\ 5 & \text{otherwise.} \end{cases}$$

Proof. Observe first that since cycles satisfy the 1-2-3 Conjecture and are regular, Corollary 2.4 implies that $\chi_{\Sigma>1}(C_\ell) \leq 5$ holds for every cycle C_ℓ .

Let $v_0, \dots, v_{\ell-1}$ denote the vertices of the cycle C_ℓ , with $v_i v_{i+1}$ being an edge for every i , $0 \leq i \leq \ell - 1$ (here and in the following, all operations over the subscripts are understood modulo ℓ). Note first that Observation 3.3 implies that no cycle C_ℓ verifies $\chi_{\Sigma>1}(C_\ell) \leq 2$. Therefore, $\chi_{\Sigma>1}(C_\ell) \geq 3$ for every cycle C_ℓ . Still according to Observation 3.3, note that no neighbour-sum-2-distinguishing 3-edge-weighting ω of C_ℓ can use weight 2. Therefore, a neighbour-sum-2-distinguishing 3-edge-weighting of C_ℓ can only use weights 1 and 3. In such a weighting, the edge weights must follow the pattern 1, 1, 3, 3, 1, 1, 3, 3, ... along the cycle, which is possible if and only if $\ell \equiv 0 \pmod{4}$.

We now prove that $\chi_{\Sigma>1}(C_\ell) = 5$ whenever $\ell \not\equiv 0 \pmod{4}$. Assume there is a neighbour-sum-2-distinguishing 4-edge-weighting ω of some cycle C_ℓ . Note that, by Observation 3.3, any two edges uv and vw being at distance 2 must be assigned a small (1 or 2) and a big (3 or 4) weight by ω . Obviously, this cannot be completely achieved when $\ell \not\equiv 0 \pmod{4}$, while it is the case otherwise (for instance, follow the pattern 1, 1, 4, 4, 1, 1, 4, 4, ... along C_ℓ). This contradicts the existence of ω , and, thus, $\chi_{\Sigma>1}(C_\ell) = 5$ whenever $\ell \not\equiv 0 \pmod{4}$. \square

3.3. Odd multicacti

Recall that $\chi_{\Sigma>1}(G) \leq 5$ holds for every odd multicactus G , according to Theorem 3.2. In the next result, we prove that, for odd multicacti, we can even design neighbour-sum-2-distinguishing 5-edge-weightings that do not use weights 2 and 4. This supports Conjecture 5.1 we raise in the concluding section.

Observe first that connected multicacti can be defined inductively, as follows. Cycles of length at least 6 and congruent to 2 modulo 4, with edges coloured green and red alternatively, are multicacti. Consider now a multicactus G whose edges are coloured green and red, and let uv be a green edge of G . Then the graph obtained from G by identifying u and v with the end-vertices of a path of length at least 5 and congruent to 1 modulo 4, whose edges are alternatively coloured red, green, ..., red (from one end to the other), is a multicactus. This operation will be referred to as a *path attachment*. Note that, in any edge-coloured multicactus, the two ends of a green edge have the same degree.

Theorem 3.6. *Every odd multicactus admits a neighbour-sum-2-distinguishing $\{1, 3, 5\}$ -edge-weighting.*

Proof. Let G be an odd multicactus. The proof is by induction on the number of path attachments performed to construct G . If no such path attachment was made, then G is a cycle C_{4k+2} , for some $k \geq 1$, and the $\{1, 3, 5\}$ -edge-weighting obtained by applying the pattern 1, 1, 3, 3, 5, 5, 1, 1, 3, 3, 5, 5, ... cyclically is clearly a neighbour-sum-2-distinguishing $\{1, 3, 5\}$ -edge-weighting of G .

Assume now that G is not a cycle. Then G must contain a green edge uv such that u and v are joined by exactly x paths P_1, \dots, P_x , $x \geq 1$, with length at least 5 and congruent

to 1 modulo 4, and whose internal vertices have degree 2 in G . In other words, no green edge of the P_i 's was used to make a path attachment. In the extremal case where the P_i 's cover all vertices of G (*i.e.*, all path attachments were performed on only one green edge), we omit one of these paths. Then, when removing all internal vertices of the P_i 's from G , we get another connected odd multicactus G' in which both u and v have degree 2. In particular, G' cannot be reduced to a single edge.

By the induction hypothesis, G' admits a neighbour-sum-2-distinguishing $\{1, 3, 5\}$ -edge-weighting ω , which we would like to extend to the edges of the P_i 's, in order to obtain a neighbour-sum-2-distinguishing $\{1, 3, 5\}$ -edge-weighting of G . Let us denote by u' and v' the neighbours of u and v , respectively, different from v and u , respectively, in G' . When extending ω to the P_i 's, we have to make sure that:

1. $\sigma(u)$ and $\sigma(u')$ still differ by at least 2;
2. $\sigma(v)$ and $\sigma(v')$ still differ by at least 2;
3. both $\sigma(u)$ and $\sigma(v)$ differ by at least 2 from the incident sums of their x neighbours in the P_i 's;
4. $\sigma(u)$ and $\sigma(v)$ still differ by at least 2.

In order to respect the fourth condition above, we will edge-weight the P_i 's in such a way that $\sigma(u)$ and $\sigma(v)$ are altered the same way, *i.e.*, by a same integer α . To that aim, we will $\{1, 3, 5\}$ -edge-weight every P_i in such a way that its two end-edges are assigned the same weight α . In what follows, the main weighting scheme we use to extend ω to the consecutive edges of some P_i is $\alpha, \beta, \beta, \gamma, \gamma, \beta, \beta, \dots, \beta, \beta, \gamma, \alpha$ (or reversely), where $\alpha, \beta, \gamma \in \{1, 3, 5\}$ and $\alpha \neq \beta$ and $\beta \neq \gamma$. Note that, by that edge-weighting, for any such α, β, γ no two adjacent inner vertices have their sums differing by strictly less than 2 (recall Observation 3.3).

We consider the worst-case scenario where the number x of P_i 's is exactly 1. This is indeed the worst case as, as x gets larger, we get more and more ways to weight the P_i 's, and thus more and more ways to extend ω correctly to G .

If $\omega(uv) = 1$, then we set $\omega(uv) = 3$. If no conflict arises, then we proceed with the next step. Otherwise, we proceed as follows. If there exists $\alpha \in \{1, 3, 5\}$ such that $\sigma(u) + \alpha$ and $\sigma(u')$ differ by at least 2, and similarly for $\sigma(v) + \alpha$ and $\sigma(v')$, then we extend ω as follows. Note first that, because $\omega(uv) = 3$, we have, say, $\sigma(u) \geq 4$ and $\sigma(v) \geq 6$ ($\sigma(u) < \sigma(v)$) by the induction hypothesis. We then extend ω to P_1 in the following way; in any case, it can be checked that the resulting edge-weighting of G is correct:

- If $\alpha = 1$, then we assign weights $1, 3, 3, 1, 1, 3, 3, \dots, 3, 3, 1, 1$ (*i.e.*, $\alpha = 1$, $\beta = 3$, $\gamma = 1$) to the edges of P_1 as going from v to u .
- If $\alpha = 3$, then we assign weights $3, 1, 1, 3, 3, 1, 1, \dots, 1, 1, 3, 3$ (*i.e.*, $\alpha = 3$, $\beta = 1$, $\gamma = 3$) as going from u to v .
- If $\alpha = 5$, then we change $\omega(uv)$ from 3 to 5 and α from 5 to 3, so that the previous case now applies.

Now, if no such α exists, then we set $\omega(uv) = 5$. In the present situation, $\alpha = 5$ now has the desired properties. Similar extension arguments then apply.

So now assume that $\omega(uv) \in \{3, 5\}$ and there is currently no sum conflict. Again, if there exists $\alpha \in \{1, 3, 5\}$ with the desired properties, then we are done. Otherwise,

note that no conflict may arise when setting $\omega(uv) = 1$ (as otherwise there would exist an α when $\omega(uv) \in \{3, 5\}$), and $\alpha = 1$ is as required. Since $\omega(uv) = 1$, we have, say, $2 \leq \sigma(u) \leq 4$ and $4 \leq \sigma(v) \leq 6$ ($\sigma(u) < \sigma(v)$). We then extend ω to the edges of P_1 as follows:

- If $\sigma(u) = 2$ and $\sigma(v) \in \{4, 6\}$, then we assign weights $1, 5, 5, 1, 1, 5, 5, \dots, 5, 5, 1, 1$ (*i.e.*, $\alpha = 1, \beta = 5, \gamma = 1$) to the edges of P_1 as going from u to v .
- If $\sigma(u) = 4$ and $\sigma(v) = 6$, then we assign weights $1, 3, 3, 1, 1, 3, 3, \dots, 3, 3, 1, 1$ (*i.e.*, $\alpha = 1, \beta = 3, \gamma = 1$) as going from v to u .

In all cases, it can be checked that no sum conflict arises in G , and thus we get a neighbour-sum-2-distinguishing $\{1, 3, 5\}$ -edge-weighting. This completes the proof. \square

4. Algorithmic aspects

In this section, we consider the hardness of determining the value of $\chi_{\Sigma>1}(G)$ for a given graph G . We first prove, in Subsection 4.1, that deciding whether $\chi_{\Sigma>1}(G) \leq 2$ holds for a given graph G is NP-complete, even when restricted to 3-degenerate planar bipartite graphs. In Subsection 4.2, we prove that, although $\chi_{\Sigma>1}(T)$ can take any value in $\{1, 2, 3, 4\}$ for a given tree T (recall Corollary 3.1), deciding the exact value of $\chi_{\Sigma>1}(T)$ can be done in polynomial time.

It is worth recalling that determining the value of $\chi_{\Sigma}(G)$ for a given graph G is NP-complete in general (Dudek and Wajc [1]), but can be done in polynomial time when restricted to bipartite graphs (Thomassen, Wu and Yang [7]). Hence our result in Subsection 4.1 shows another difference between the parameters χ_{Σ} and $\chi_{\Sigma>1}$.

4.1. General case

Before proceeding with the proof of the main result of this subsection, we first introduce gadgets that we will use to force some weights to be used by any neighbour-sum-2-distinguishing 2-edge-weighting. Each of these gadgets will have a *root vertex* of degree 1 being incident to a *root edge*. We here relax the notion of neighbour-sum-2-distinguishing 2-edge-weighting around the root; that is, we allow a neighbour-sum-2-distinguishing 2-edge-weighting to have adjacent incident sums differing by less than 2, but the incident sum of the root has to be involved in such a conflict. This is because our gadgets will be attached to other graphs via the root, so, in the properties we point out below, the incident sum of the root should not be regarded as fixed.

The gadgets we will construct are called (α, S) -gadgets, for some given $\alpha \in \{1, 2\}$ and $S \subset \mathbb{N}^*$ ($S \neq \emptyset$). Every such gadget G will satisfy the two following properties:

1. the root edge of G is necessarily weighted α by any neighbour-sum-2-distinguishing 2-edge-weighting of G , and
2. S is the set of all numbers s such that there exists a neighbour-sum-2-distinguishing 2-edge-weighting of G where the (unique) neighbour of the root has incident sum s .

These gadgets will be used as follows. Let H be a graph and v be a vertex of H . Add to H an (α, S) -gadget G (for some α and S), and identify v with the root of G . Then, in any neighbour-sum-2-distinguishing 2-edge-weighting of H , v will necessarily receive weight α from the root edge of G , and v will be adjacent to a vertex whose incident sum belongs to S . This mechanism can be used both to force particular edge weights to appear around v , and to ensure that $\sigma(v)$ does not belong to a particular set (in particular when $|S| = 1$).

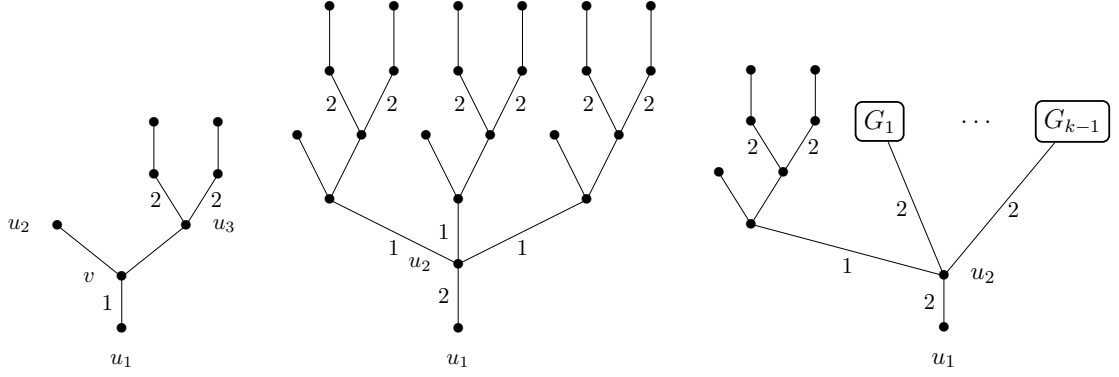


Figure 1: The $(1, \{3, 4\})$ -gadget with root u_1 (left), the $(2, \{5\})$ -gadget with root u_1 (middle), and the $(2, \{2k+1\})$ -gadget with root u_1 (right).

We now introduce the gadgets we will use (see Figure 1 for an illustration). Consider first a path $G := u_1 u_2 u_3$ of length 2. We claim that G is a $(2, \{3, 4\})$ -gadget with root u_1 . Indeed, in any neighbour-sum-2-distinguishing 2-edge-weighting ω of G (with the relaxation mentioned above), we have $\omega(u_1 u_2) = 2$, while $\omega(u_2 u_3)$ can have value either 1 or 2, in which cases we get $\sigma(u_2) = 3$ and $\sigma(u_2) = 4$, respectively.

Now consider a claw G with vertices v, u_1, u_2, u_3 , where the u_i 's are the leaves. Add two $(2, \{3, 4\})$ -gadgets G_1 and G_2 to G , and identify u_3 and the roots of G_1 and G_2 . We claim that G is a $(1, \{3, 4\})$ -gadget with root u_1 . In any neighbour-sum-2-distinguishing 2-edge-weighting ω of G , the vertex u_3 is incident to at least two edges with weight 2 (because of the gadgets G_1 and G_2), so that $\sigma(u_3) = 5$ (if $\omega(u_3 v) = 1$) or $\sigma(u_3) = 6$ (otherwise). In both cases, we necessarily have $\omega(v u_2) = \omega(v u_1) = 1$, so that $\sigma(v)$ and $\sigma(u_3)$ differ by at least 2. We thus get $\sigma(v) = 3$ in the first case, and $\sigma(v) = 4$ in the second case.

We now describe how to obtain $(2, S)$ -gadgets with $S := \{2k+1\}$ for any $k \geq 2$. We first build a $(2, \{5\})$ -gadget as follows. Start from $G := u_1 u_2$ being the path of length 1, then add three $(1, \{3, 4\})$ -gadgets G_1, G_2, G_3 to G , and identify u_2 and the roots of G_1, G_2 and G_3 . We claim that G is a $(2, \{5\})$ -gadget with root u_1 . In any neighbour-sum-2-distinguishing 2-edge-weighting ω of G , the vertex u_2 is incident to at least three edges weighted 1, namely the root edges of the G_i 's. Now, if $\omega(u_1 u_2) = 1$, then $\sigma(u_2) = 4$, which creates sum conflicts with vertices from the G_i 's. So we necessarily have $\omega(u_1 u_2) = 2$, in which case $\sigma(u_2) = 5$, which is fine since the G_i 's are $(1, \{3, 4\})$ -gadgets.

We now turn to the general case. Let $2k+1 \geq 7$, and assume that we have constructed $(2, S')$ -gadgets with $S' := \{2k'+1\}$ for every $k', 2 \leq k' < k$. Start from $G := u_1 u_2$ being the path of length 1. Add $k-1$ $(2, \{2k-1\})$ -gadgets G_1, \dots, G_{k-1} to G , as well as one $(1, \{3, 4\})$ -gadget G_0 , and identify u_2 and the roots of G_1, \dots, G_{k-1} and G_0 . We claim that G is a $(2, \{2k+1\})$ -gadget with root u_1 . In any neighbour-sum-2-distinguishing 2-edge-weighting ω of G , the G_i 's force $\sigma(u_2)$ to have value at least $2k-1$. Depending on whether $\omega(u_1 u_2) = 1$ or $\omega(u_1 u_2) = 2$, we thus have $\sigma(u_2) = 2k$ or $\sigma(u_2) = 2k+1$, respectively. In the first case, we get sum conflicts between u_2 and its neighbours in G_1, \dots, G_{k-1} since their incident sums differ by 1. Therefore, we necessarily have $\omega(u_1 u_2) = 2$, so that $\sigma(u_2) = 2k+1$, which produces no sum conflict in G .

Analogous $(1, S)$ -gadgets with $S := \{2k+1\}$, $k \geq 2$, will also be needed. A $(1, \{5\})$ -gadget can be obtained as follows. Start from $G := u_1 u_2$ being the path of length 1, add two $(2, \{7\})$ -gadgets G_1 and G_2 to G , and identify u_2 and the roots of G_1 and G_2 . In any neighbour-sum-2-distinguishing 2-edge-weighting ω of G , the G_i 's force u_2 to be incident to at least two edges with weight 2. So we have $\sigma(u_2) = 5$ or $\sigma(u_2) = 6$ depending on whether

$\omega(u_1u_2) = 1$ or $\omega(u_1u_2) = 2$, respectively. In the second case, however, we get sum conflicts between u_2 and its neighbours in G_1 and G_2 . We thus necessarily have $\omega(u_1u_2) = 1$ and $\sigma(u_2) = 5$, which produces no sum conflict in G .

Now let $2k + 1 \geq 7$, and assume that we have constructed $(1, S')$ -gadgets with $S' := \{2k' + 1\}$ for every k' , $2 \leq k' < k$. Start from $G := u_1u_2$ being the path of length 1, add k $(2, \{2k + 3\})$ -gadgets G_1, \dots, G_k to G , and identify u_2 and the roots of G_1, \dots, G_k . We claim that G is a $(1, \{2k + 1\})$ -gadget. In any neighbour-sum-2-distinguishing 2-edge-weighting ω of G , the G_i 's force $\sigma(u_2)$ to be at least $2k$. Depending on whether $\omega(u_1u_2) = 1$ or $\omega(u_1u_2) = 2$, we thus have $\sigma(u_2) = 2k + 1$ or $\sigma(u_2) = 2k + 2$, respectively. In the second case, we get sum conflicts between u_2 and its neighbours in G_1, \dots, G_k , since their incident sums differ by 1. Therefore, we necessarily have $\omega(u_1u_2) = 1$ and $\sigma(u_2) = 2k + 1$, which produces no sum conflict in G .

Note that all the above-constructed gadgets are trees. With all these gadgets in hand, we now prove the main result of this section.

Theorem 4.1. *For a given 3-degenerate planar bipartite graph G , deciding whether $\chi_{\Sigma > 1}(G) \leq 2$ holds is NP-complete.*

Proof. Since the problem is obviously in NP, we proceed with the proof of its NP-hardness. The proof is by reduction from 1-IN-3 SAT. From a formula F , we construct a graph G such that F is satisfiable in a 1-in-3 way if and only if G admits a neighbour-sum-2-distinguishing 2-edge-weighting. Since the MONOTONE version of 1-IN-3 SAT remains NP-complete (see e.g. [5]), we may assume that F has no negated variables. Also, we may assume that all clauses of F have three distinct variables, as otherwise F could be simplified. That is:

- if F has a clause $(x_{i_1} \vee x_{i_1} \vee x_{i_1})$, then F is not satisfiable in a 1-in-3 way;
- if F has a clause $(x_{i_1} \vee x_{i_1} \vee x_{i_2})$, then x_{i_2} and x_{i_1} are forced to *true* and *false*, respectively, by any truth assignment making F satisfied in a 1-in-3 way.

We denote by x_1, \dots, x_n the variables of F , and by C_1, \dots, C_m its clauses. The construction of G , which is clearly achieved in polynomial time, is as follows. We start by adding *variable gadgets* in the following way. For each variable x_i of F , we add to G a star V_i with root v_i and $2k_i$ leaves $u_{i,1}, \dots, u_{i,2k_i}$, where $2k_i \geq \max\{10, n_i\}$ is any even integer, and n_i is the number of clauses of F that contain x_i . Next we add $(1, \{4k_i + 1\})$ -, $(2, \{4k_i + 3\})$ -, $(2, \{4k_i + 5\})$ -, \dots , $(2, \{6k_i - 3\})$ - and $(2, \{6k_i - 1\})$ -gadgets G_1, \dots, G_{k_i} to G , and identify v_i and the roots of G_1, \dots, G_{k_i} . To avoid any ambiguity, let us emphasize that, for instance, when $k_i = 5$, the five added gadgets are one $(1, \{21\})$ -gadget, one $(2, \{23\})$ -gadget, one $(2, \{25\})$ -gadget and two $(2, \{27\})$ -gadgets.

Because of the G_i 's, in any neighbour-sum-2-distinguishing 2-edge-weighting ω of V_i , the value of $\sigma(v_i)$ lies between $4k_i - 1$ (when all $v_iu_{i,j}$'s are assigned weight 1) and $6k_i - 1$ (when all $v_iu_{i,j}$'s are assigned weight 2). Furthermore, $4k_i - 1$ and $6k_i - 1$ are both odd. Moreover, we cannot have $\sigma(v_i) \in \{4k_i, \dots, 6k_i - 2\}$ as otherwise there would be a sum conflict involving v_i and one of its neighbours in the G_i 's. Therefore, either all $\omega(v_iu_{i,j})$'s are equal to 1, or all $\omega(v_iu_{i,j})$'s are equal to 2. In what follows, we call the vertices $u_{i,1}, \dots, u_{i,2k_i}$ the *output vertices* of V_i , and the edges $v_iu_{i,1}, \dots, v_iu_{i,2k_i}$ the *output edges* of V_i .

We now modify G by considering the clauses of F . For each clause $C_j = (x_{i_1} \vee x_{i_2} \vee x_{i_3})$ of F , we add a *clause vertex* c_j to G . For each $V_{i_1}, V_{i_2}, V_{i_3}$, we then select one output vertex still having degree 1, and identify c_j and the three selected output vertices. Finally, we add

a $(2, \{7\})$ -gadget G_1 to G , as well as a $(2, \{11\})$ -gadget G_2 , and identify c_j and the roots of G_1 and G_2 . In any neighbour-sum-2-distinguishing 2-edge-weighting ω of G , $\sigma(c_j)$ has thus value at least 4 (because of G_1 and G_2), and ranges in $\{7, \dots, 10\}$. However, $\sigma(c_j)$ cannot take any value among $\{7, 8, 10\}$ because of G_1 and G_2 . So we necessarily have $\sigma(c_j) = 9$, which occurs only if exactly one of the three output edges originating from $V_{i_1}, V_{i_2}, V_{i_3}$ is assigned weight 1.

It can be checked that no unexpected sum conflicts (that is, different from those listed above) can arise, in particular thanks to our choice of the $2k_i$'s. We now claim that we have the desired equivalence. This directly follows from the following arguments:

- Assigning weight 1 (resp. 2) to an output edge $v_i c_j$ simulates the fact that variable x_i brings truth value *true* (resp. *false*) to C_j .
- Following that equivalence, the fact that, for any V_i , all output edges of V_i must be weighted 1 (resp. 2) simulates the fact that setting x_i to *true* (resp. *false*) by some truth assignment brings the same truth value to every clause containing x_i .
- The fact that, for every clause vertex c_j , exactly one incident output edge must be assigned weight 1 simulates the fact that a clause of F is considered satisfied if and only if it includes exactly one *true* variable.

To complete the proof of the theorem, it suffices to observe the following:

- The PLANAR version of MONOTONE 1-IN-3 SAT remains NP-complete (see [5]), so we may assume that F is a planar formula. Since every gadget is a tree, the construction above then yields a planar G .
- Since every gadget is a tree, the graph G is 3-degenerate.
- The only cycles in G are those of the subgraph induced by the v_i 's and the c_j 's. Since this subgraph is bipartite, so is G .

This concludes the proof. □

4.2. Tree case

In this section, we prove that the counterpart of Theorem 4.1 for trees is not true. That is, we prove that determining the value of $\chi_{\Sigma>1}(T)$ for a given tree T can be done in polynomial time. Recall that for a tree T , we always have $\chi_{\Sigma>1}(T) \leq 4$ (according to Corollary 3.1), while $\chi_{\Sigma>1}(T) = 1$ if and only if, for every two adjacent vertices u and v of T , the values of $d(u)$ and $d(v)$ differ by at least 2.

Theorem 4.2. *For a given tree T , determining $\chi_{\Sigma>1}(T)$ can be done in polynomial time.*

Proof. For any fixed $k \in \{1, 2, 3, 4\}$, we introduce below an algorithm that checks in polynomial time whether T admits a neighbour-sum-2-distinguishing k -edge-weighting. So, to determine $\chi_{\Sigma>1}(T)$, we can essentially run this algorithm successively with $k = 1, 2, 3, 4$. The first value of k for which the algorithm answers positively is the value of $\chi_{\Sigma>1}(T)$.

Designate a node r of T as being its *root*. This defines a root-to-leaf orientation of T in the usual way, where every non-root node v has a *parent*, and every non-leaf node v has *children*. By the *descendants* of v , we refer to the nodes of T for which we find v when iterating the parent relationship.

The subtree T_v of T rooted at v is the subtree whose nodes are v and all its descendants. This subtree T_v can itself be decomposed into several subtrees, in the following way. Assume that v has $d \geq 1$ descendants u_1, \dots, u_d , ordered following an arbitrary order (supposed to be fixed throughout the proof). Then T_v can be edge-decomposed into d subtrees $T_{v,1}, \dots, T_{v,d}$ being $T_{u_1} + vu_1, \dots, T_{u_d} + vu_d$, respectively, whose root, v , has degree precisely 1. Trees with this property are called *shrubs* throughout. Conversely, T_v is obtained by identifying the roots of the shrubs $T_{v,1}, \dots, T_{v,d}$. For every shrub, we call the edge incident to the root the *root edge*. The non-root end of the root edge is called the *subroot*.

We are now ready to describe our algorithm for deciding whether T admits a neighbour-sum-2-distinguishing k -edge-weighting. The rough ideas are the following. The tree T can be seen as a union of $d := d(r)$ shrubs S_1, \dots, S_d whose roots were identified, resulting in r . A neighbour-sum-2-distinguishing k -edge-weighting of T is thus essentially the union of (relaxed, see below) neighbour-sum-2-distinguishing k -edge-weightings of the d shrubs attached to r , with the additional property that the resulting $\sigma(r)$ does not create any sum conflict. Therefore, in order to construct a neighbour-sum-2-distinguishing k -edge-weighting of T , it suffices to find convenient neighbour-sum-2-distinguishing k -edge-weightings of S_1, \dots, S_d that can be “glued”. So we need to know, for each shrub S_i and for every $\alpha \in \{1, \dots, k\}$, whether S_i admits a neighbour-sum-2-distinguishing k -edge-weighting where the root edge is assigned colour α , and, for such an edge-weighting of S_i , which possible incident sums can be obtained for the subroot.

More formally, for a shrub S with root v' and subroot v , we want to compute, for every weight $\alpha \in \{1, \dots, k\}$, the set $X_\alpha(v)$ of possible values of $\sigma(v)$ by a neighbour-sum-2-distinguishing k -edge-weighting of S assigning weight α to $v'v$. Note that a shrub might be a single edge, and may thus admit no neighbour-sum-2-distinguishing k -edge-weighting. In that special case, we relax the notion of neighbour-sum-2-distinguishing k -edge-weighting, and allow the root and the subroot to have the same incident sums.

Assume v has d children u_1, \dots, u_d , $d \geq 0$, and let S_1, \dots, S_d denote the d shrubs attached to v in S . We claim that each $X_\alpha(v)$ can be computed in polynomial time from the sets

$$X_1(u_1), \dots, X_k(u_1), X_1(u_2), \dots, X_k(u_2), \dots, X_1(u_d), \dots, X_k(u_d),$$

computed by induction for the shrubs S_1, \dots, S_d . So, in a way, the sets $X_1(v), \dots, X_k(v)$ can be computed from smaller shrubs, and deduced successively towards the subroot of S . We prove this below.

The base case is when S is a single edge, that is, v has no child. If the edge $v'v$ is assigned any weight $\alpha \in \{1, \dots, k\}$ by a neighbour-sum-2-distinguishing k -edge-weighting, then $\sigma(v) = \alpha$. So, for such a shrub S , we have $X_\alpha(v) = \{\alpha\}$ for every $\alpha \in \{1, \dots, k\}$.

Suppose now that v has $d \geq 1$ children u_1, \dots, u_d , and, for each shrub S_i , $1 \leq i \leq d$, attached to v , and every $\alpha \in \{1, \dots, k\}$, the set $X_\alpha(u_i)$ has been computed by induction. We now want to compute the sets $X_1(v), \dots, X_k(v)$. Since $d(v) = d + 1$, by any k -edge-weighting of S , the sum $\sigma(v)$ can take up to $kd + k - d$ values, namely those among $\{d + 1, \dots, kd + k\}$. We repeatedly fix one of those sums x , and we determine whether x can be added to some of the sets $X_1(v), \dots, X_k(v)$.

Assume we want to determine whether x has to be added to $X_\alpha(v)$, where α is any value in $\{1, \dots, k\}$. Successively consider all partitions $x_1 + 2x_2 + \dots + kx_k$ of x into $x_1 + \dots + x_k = d + 1$ values among $\{1, \dots, k\}$ only. Recall that we are focusing on computing $X_\alpha(v)$, so if $x_\alpha = 0$, then we can consider the next partition of x . Since x is linear in $|V(T)|$ and $k \leq 4$ is fixed, the number of such partitions to consider is polynomial in $|V(T)|$. Essentially, we have $x_1, \dots, x_k \leq |V(T)|$, meaning that the number of such

partitions is roughly $|V(T)|^{k-1}$. We now want to know if there is a neighbour-sum-2-distinguishing k -edge-weighting of S where x_i edges incident to v are assigned weight i , for every $i \in \{1, \dots, k\}$. If such an edge-weighting exists, then for any S_i to which the weight β is assigned to the root edge, $X_\beta(u_i)$ contains some value not in $\{\alpha - 1, \alpha, \alpha + 1\}$.

Since we are focusing on $X_\alpha(v)$, one of the x_α weights α around v will be assigned to $v'v$. This leaves us with d other weights to assign bijectively to the vu_i 's, with the constraint that if we assign a weight β to vu_i , then $X_\beta(u_i)$ should contain a value not among $\{\alpha - 1, \alpha, \alpha + 1\}$. If β can indeed be assigned to vu_i safely, then we call this a *valid assignment*. To find a correct assignation (if any exists), we build a compatibility bipartite graph C of the valid assignments, as follows. In one side of the bipartition of C , we put d vertices corresponding to the d weights we want to assign. In the other side, we put d vertices corresponding to the edges vu_1, \dots, vu_d of S . We then add an edge joining two vertices of C if assigning the corresponding weight to the corresponding edge of S is valid. Now, finding a satisfying assignment of the d weights to the root edges of the S_i 's is equivalent to finding a perfect matching in C , which is known to be doable in polynomial time. If there indeed exists such a perfect matching of C , then we add x to $X_\alpha(v)$.

We now go back to T , with the root r having d children u_1, \dots, u_d . For each S_i of the d shrubs S_1, \dots, S_d rooted at v , we can compute the sets $X_1(u_i), \dots, X_k(u_i)$ as explained above. These sets memorize, in a compact way, all possible ways, in terms of incident sums and weights assigned to the root edges, to k -edge-weight the S_i 's in a neighbour-sum-2-distinguishing way. Now, again, we can consider every potential incident sum x as $\sigma(r)$, every potential way to partition x into d integers among $\{1, \dots, k\}$, and, building the compatibility bipartite graph as above, find, if it exists, a valid way to bijectively assign the d weights to the d root edges vu_1, \dots, vu_d . If a valid assignment for a partition of some x exists, then T admits a neighbour-sum-2-distinguishing k -edge-weighting. Otherwise, it does not.

Concerning the complexity aspect, determining $X_\alpha(v)$ for a shrub of T with subroot v can be done in polynomial time. Recall that $k \leq 4$ is constant. The number of possible sums x as $\sigma(v)$ to consider is at most $k|V(T)|$. For each of these values of x , we consider up to $|V(T)|^{k-1}$ partitions into 1's, 2's, \dots , and k 's. Deciding whether there is a valid assignment for one of those partitions can be done in polynomial time, using for instance Edmonds' Blossom Algorithm for computing maximum matchings [2]. The procedure above is almost the same when r is considered. By all these arguments, the whole procedure can be achieved in polynomial time. \square

5. Conclusion

In this work, we have investigated the consequences on the 1-2-3 Conjecture of requiring adjacent vertices to be distinguishable in a stronger way, namely by asking their incident sums to differ by at least 2. We have addressed Conjecture 1.1, to which we did not manage to come up with any counterexample, as an equivalent of the 1-2-3 Conjecture in this context. As a main evidence that our conjecture might be true, we have pointed out some connections between the 1-2-3 Conjecture and Conjecture 1.1, and proved the later one for nice bipartite graphs.

Several aspects related to Conjecture 1.1 remain unclear to us, and could thus be subject to further work. First, we do not fully understand how necessary the weights 2 and 4 are for our conjecture. In particular, most graphs for which we have proved Conjecture 1.1 actually admit neighbour-sum-2-distinguishing $\{1, 3, 5\}$ -edge-weightings (recall, in particular, Theorem 3.6). This supports the following refinement of Conjecture 1.1.

Conjecture 5.1. *Every nice graph admits a neighbour-sum-2-distinguishing $\{1, 3, 5\}$ -edge-weighting.*

In the context of neighbour-sum-2-distinguishing edge-weightings, Conjecture 5.1 might actually be an equivalent to the 1-2-3 Conjecture more natural than Conjecture 1.1. Indeed, in the 1-2-3 Conjecture we aim at getting incident sums differing by at least 1 by using three successive weights $\alpha - 1, \alpha, \alpha + 1$ differing by 1. In Conjecture 5.1, we aim at getting incident sums differing by at least 2 by using three “successive” weights $\alpha - 2, \alpha, \alpha + 2$ differing by 2. Following this reasoning, perhaps, in general, the following conjecture might be the right direction to consider.

Conjecture 5.2. *Let $d \geq 1$ be an integer. Every nice graph admits a neighbour-sum- d -distinguishing $\{1, d + 1, 2d + 1\}$ -edge-weighting.*

There are intriguing examples, though, such as nice paths P_ℓ of length congruent to 1 modulo 4 (for which $\chi_{\Sigma > 1}(P_\ell) = 4$, recall Theorem 3.4), showing that, for neighbour-sum-2-distinguishing edge-weightings, the weights 2 and 4 are sometimes worth using to get an optimal edge-weighting. More generally, our NP-hardness reduction in the proof of Theorem 4.1 shows that there exist complex examples for which an optimal edge-weighting uses weights 1 and 2 only. We believe this could be an interesting aspect to study further.

Although we have proved Conjecture 1.1 for nice bipartite graphs (Theorem 3.2), we have not proved refined Conjecture 5.1 for all such graphs. Another interesting line of research could thus be to generalize Theorem 3.6 to all nice bipartite graphs, which would be a first step towards Conjecture 5.1.

More directions for future works on neighbour-sum-2-distinguishing edge-weightings are also worth mentioning. Notably, we did not manage to improve the bounds given in Section 2 for many classes of graphs. Generally speaking, it does not seem obvious to us how to improve the bound in Corollary 2.2, and this would surely require new dedicated tools. Concerning particular classes of graphs, let us mention the case of subcubic graphs. Although we know that cubic graphs comply with Conjecture 1.1, and even Conjecture 5.1 (recall Corollary 2.4), we did not manage to prove that nice subcubic graphs, in general, also do. We believe this would be an appealing first case to consider towards proving Conjecture 1.1 for 3-chromatic graphs, for which the 1-2-3 Conjecture holds.

More generally, it would be interesting to consider Conjecture 5.2 above. Many of the arguments and techniques used in this work actually generalize to neighbour-sum-distinguishing d -edge-weightings. For values of d larger than 2, it is likely that more intriguing phenomenon arise.

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